

Numerical simulation with MPS

(89)

- Evaluation of expectation values is of the form

$$\begin{array}{c} \square - \square - \dots - \square - \square \\ | \quad | \quad \quad \quad | \quad | \\ \square - \square - \dots - \square - \square \end{array} = (e | E \cdot \dots \cdot E \cdot | r)$$

\uparrow
 $D^2 \times D^2$ - matrices.

Start from left:

$$(1 \times D^2) \cdot (D^2 \times D^2) \text{ - multiplication: } D^4 \text{ steps}$$

\Rightarrow total $\sim ND^4$ steps \Rightarrow efficient in N (as long as $D = \text{poly}(N)$)

(Note: D^6 for PBC)

\Rightarrow Efficient computation of exp. values etc.!

- Given local Hamiltonian $H = \sum h_i$, can we find the best MPS approx. to its ground state, i.e.

MPS $|\psi\rangle$ w/ some D s.t. $\langle \psi | H | \psi \rangle$ minimal?

(H gapped \rightarrow area law \rightarrow will be a good approx. of GS.)

Will present basic ideas first. Several omissions / "bad practice"
- will be fixed later!

Consider MPS $|\psi\rangle \equiv \sum_k (A^{[1]k_1} A^{[2]k_2} \dots A^{[N]k_N}) |k_1 - k_N\rangle$

(90)

$|\psi\rangle$ is a linear function $|\psi(A^{[s]})\rangle$ in each $A^{[s]}$ individually.

Then, $\langle \psi(A^{[s]}) | H | \psi(A^{[s]}) \rangle$ is quadratic in $A^{[s]}$,

$$\langle \psi(A^{[s]}) | H | \psi(A^{[s]}) \rangle =: \vec{A}^{[s]\dagger} \cdot \Pi \cdot \vec{A}^{[s]}$$

↑
"vectorized version" of $A^{[s]}$.

Can Π be computed efficiently? — Yes (cf. later, but
using a matrix series would be (in principle) efficient.)
↑
don't do this...

Similarly: $\langle \psi(A^{[s]}) | \psi(A^{[s]}) \rangle = \vec{A}^{[s]\dagger} \cdot N \cdot \vec{A}^{[s]}$

Can we minimize $E(A^{[s]}) = \frac{\vec{A}^{[s]\dagger} \Pi \vec{A}^{[s]}}{\vec{A}^{[s]\dagger} N \vec{A}^{[s]}}$

to find the optimal $A^{[s]}$?

→ Maps to an eigenvalue problem:

$$\min_{\vec{x}} \frac{\vec{x}^T M \vec{x}}{\vec{x}^T N \vec{x}} = \min_{\vec{x}} \frac{(\frac{1}{\sqrt{N}} \vec{x})^T \frac{1}{N} M \frac{1}{\sqrt{N}} (\frac{1}{\sqrt{N}} \vec{x})}{(\frac{1}{\sqrt{N}} \vec{x})^T (\frac{1}{\sqrt{N}} \vec{x})} = \min_{\vec{y}} \frac{\vec{y}^T \frac{1}{N} M \frac{1}{N} \vec{y}}{|\vec{y}|^2} \quad (91)$$

\Rightarrow solution is smallest eigenvalue $\lambda_{\min} \left(\frac{1}{N} M \frac{1}{N} \right)$.

(Corresponds to "generalized eigenvalue problem" $M \vec{y} = \lambda N \vec{y}$.)

\Rightarrow Efficiently yields optimal $A^{(k)}$ for a fixed configuration of the other $A^{[k]}$.

o Iterative algorithm:

- o Start w. some state (random, good guess, other meth.)
- o Start at site $s=1$. Optimize $A^{[1]}$.
- o Repeat for $s=2, 3, \dots, N$. ("sweep")
- o Repeat for $s=N-1, \dots, 1$,
- o and so on ...

\rightarrow energy will always decrease

\rightarrow monitor energy, stop when suff. converged.

Core idea of the Density Matrix Renormalization Group algorithm (DMRG, White '92)

• Is this provably necessary?

→ No. Counterexamples exist!

→ But: Works extremely well in practice (w/ well-chosen initial state...)

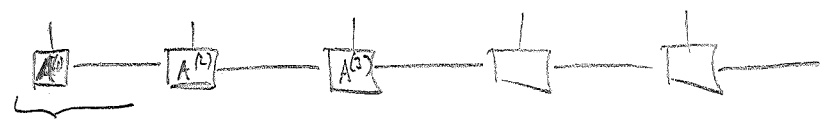
→ Provably conv. efficient algorithm recently found (very different approach & not (yet?) practical, Landau, Vazirani, Vidick '13)

Improvements / Implementation issues to keep in mind:

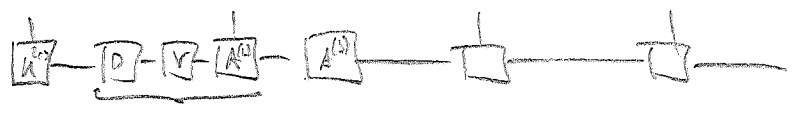
① $\frac{\vec{A} \cdot \Pi \vec{A}}{\vec{A} \cdot N \cdot \vec{A}}$ can be ill-conditioned if N has small eigenvalues.

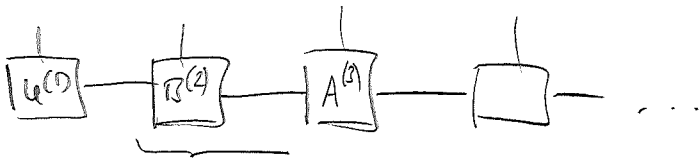
Can we resolve that / get out of N ?

→ Restrict to OBC. Then we can choose an "isometric gauge" (cf. the MPS construction from an arbitrary state in lecture 10, pg. 71-74):

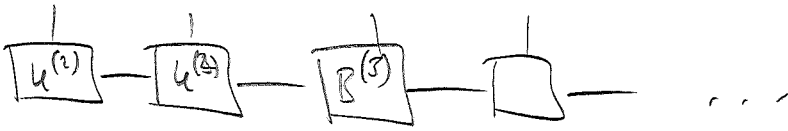
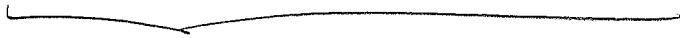
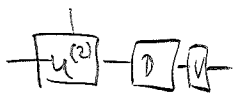


SVD:





SVD



etc.

Properties: $U^{(i)}$ isometry:

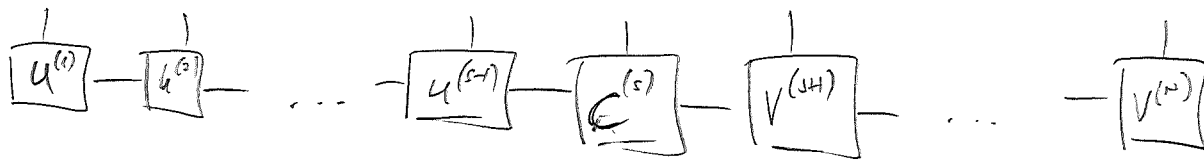
$$\begin{matrix} \boxed{u^{(i)}} & \xrightarrow{\alpha'} \\ \downarrow \\ \boxed{u^{(i)}} & \xrightarrow{\alpha} \end{matrix} = \begin{bmatrix} \alpha \\ \alpha' \end{bmatrix} \equiv \delta_{\alpha\alpha'}$$

$$\begin{matrix} \boxed{u^{(k)}} \\ \downarrow \\ \boxed{u^{(k)}} \end{matrix} = \begin{bmatrix} \\ \end{bmatrix}$$

We can do this also from the right end:

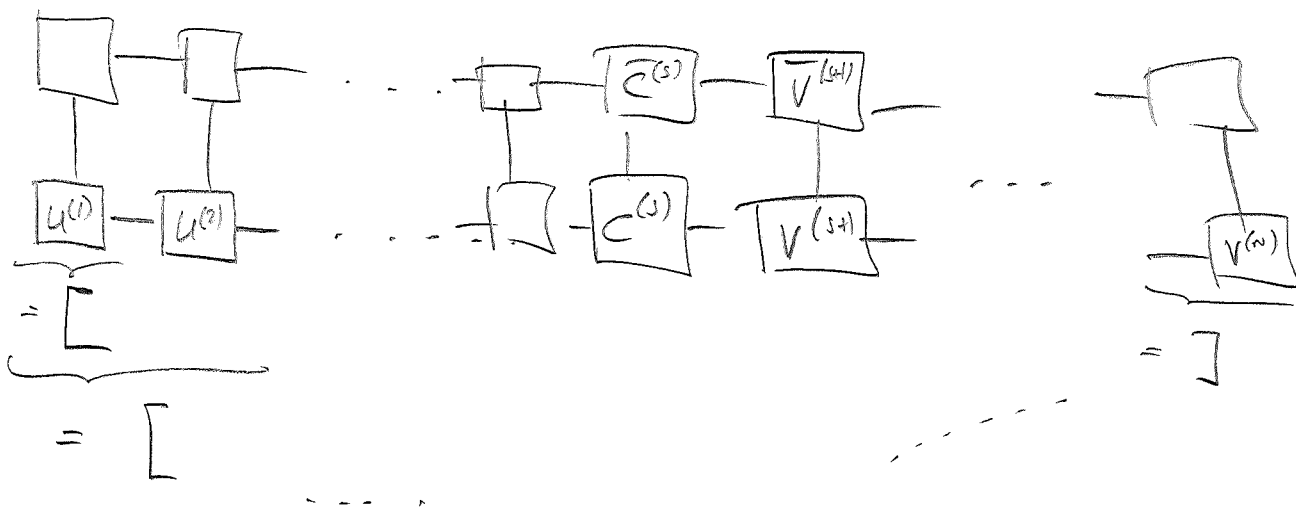
$$\begin{matrix} \boxed{v^{(n)}} \\ \downarrow \\ \boxed{v^{(n)}} \end{matrix} = \begin{bmatrix} \\ \end{bmatrix} ; \quad \begin{matrix} \boxed{v^{(k)}} \\ \downarrow \\ \boxed{v^{(k)}} \end{matrix} = \begin{bmatrix} \\ \end{bmatrix}$$

Perform gauge around site s :



(Note: $c^{(s)}$ contains same info as $A^{(s)}$ - we can as well optimize $c^{(s)}$)

Now what is $\langle \psi(c^{(s)}) | \psi(c^{(s)}) \rangle$?



$$= \left[\begin{array}{c} \bar{c}^{(s)} \\ | \\ c^{(s)} \end{array} \right] = \frac{|\vec{c}^{(s)}|^2}{|c^{(s)}|} !$$

\Rightarrow Have removed $N \rightarrow$ optimization problem

$$\min_{|c^{(s)}|=1} \frac{\langle \psi(c^{(s)}) | H | \psi(c^{(s)}) \rangle}{\langle \psi(c^{(s)}) | \psi(c^{(s)}) \rangle} = \min_{|c^{(s)}|=1} \vec{c}^{(s)T} \tilde{H} \vec{c}^{(s)}$$

is a normal eigenvalue problem (well-cond.!!)

- o Saves computational effort (no N needed, + see later)
- o When sweeping, updating gauge only requires to update one site!

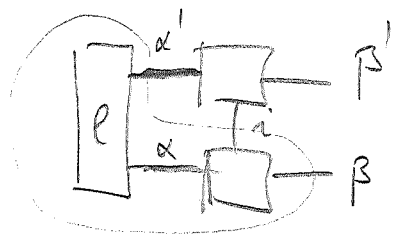
Further optimizations:

② Tensor contraction:

Head: $\langle 4|0|4 \rangle = (e| \cdot E \cdot E \dots \text{etc.}$

$$\Rightarrow (1 \times D^2) \cdot (D^2 \times D^2) = D^4,$$

But:



can be summed in different orders, and type.

$D \gg d$ (since large $D \Rightarrow$ good approx.)

Eg.: first sum α : $(D \times D) \cdot (D \times dD) \rightarrow dD^3 \text{ ops.}$

then sum α' : $(D \times dD) \cdot (dD \times D) \rightarrow dD^3 \text{ ops.}$

$$\Rightarrow dD^3 \text{ ops.} \ll D^4 \text{ ops!}$$

\Rightarrow Scaling of method is $D^3 N$!

③ Optimization in computation of Hamiltonian:

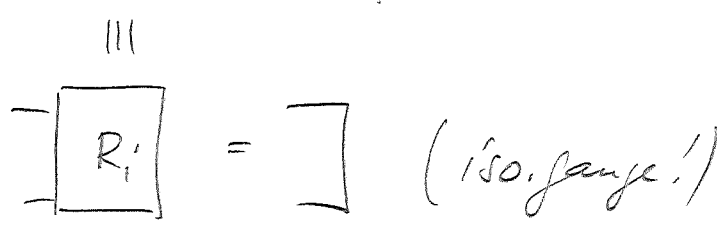
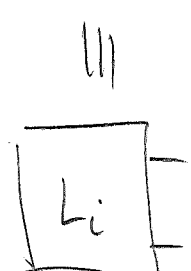
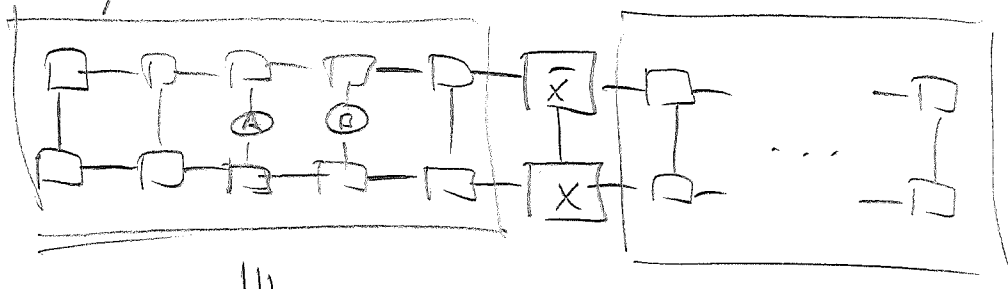
(Different approaches exist, ...)

$$\langle \psi(A^{[s]}) | H | \psi(A^{[s]}) \rangle = \vec{A}^{[s] \dagger} \Pi \vec{A}^{[s]} \Rightarrow \Pi = ?$$

Use $H = \sum h_i \Rightarrow \Pi = \sum \Pi_i$.

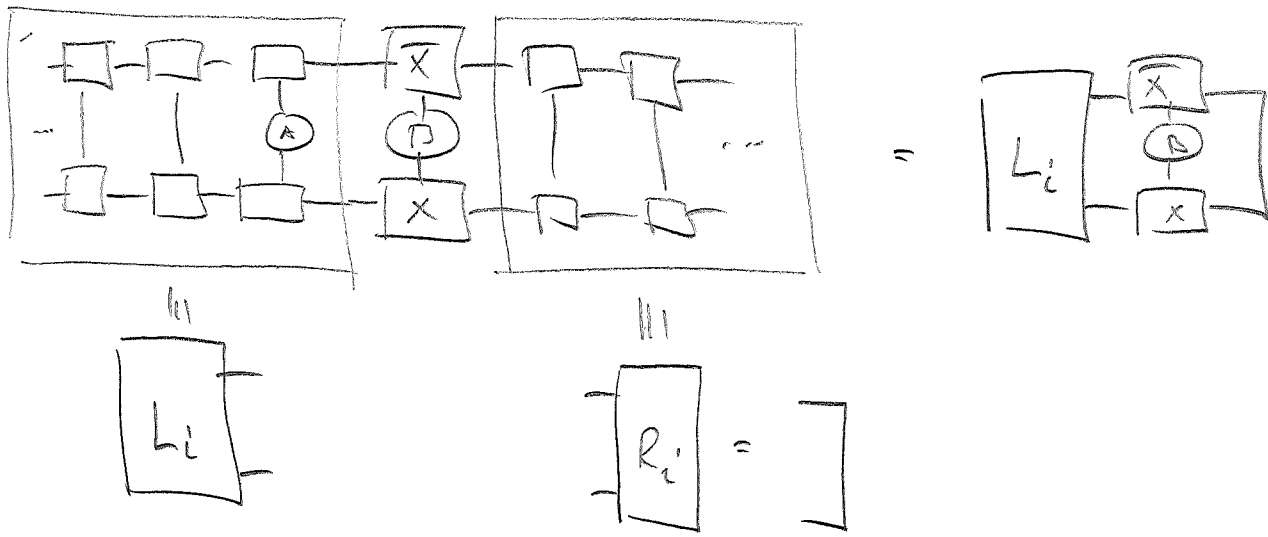
Assume wlog $h_i = A \otimes B$.

Different types of tensors: ($X \equiv A^{[s]}$)



$$= \left[\begin{array}{c} \text{---} \alpha' \\ \boxed{L_i} \\ \text{---} \alpha \end{array} \right] \begin{array}{c} \boxed{X} \\ | \\ \boxed{L_i} \\ \boxed{X} \end{array} \left[\begin{array}{c} \text{---} \beta \\ \text{---} \beta \end{array} \right] \Rightarrow \Pi_i = \underset{i}{\mathbb{1}} \otimes \underset{\alpha, \alpha'}{L_i} \otimes \underset{\beta}{\mathbb{1}}$$

... similar for A, B right of s.



$$\Rightarrow \Pi_i = B \otimes L_i \otimes I$$

$$\Pi = \sum \Pi_i \text{ over all Ham. terms.}$$

Optimizations for Π_i

- The whole influence of all Ham. terms left of a position s can be summed up in

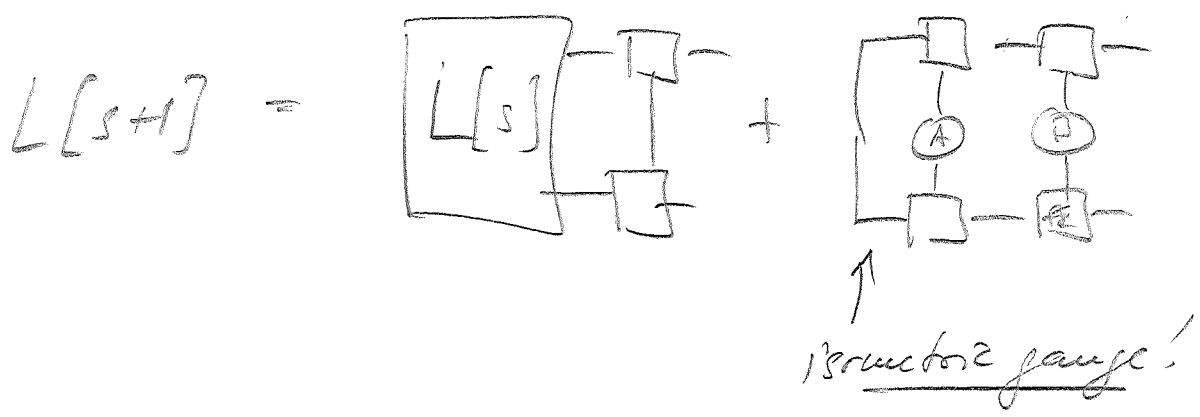
$$L[s] = \sum_{iH < s} \text{Diagram}$$

~~When moving $s \rightarrow s+1$, $L[s+1]$ can be computed from~~

~~$L[s]$~~

Ham. terms

• When moving from $s \rightarrow s+1$, $L[s+1]$ can be computed from $L[s]$ as



• By storing all $L[k]$ for $k \leq s$, we already know $L[s]$ when moving $s \rightarrow s-1$.

• Correspondingly for right!

⇒ Effort to determine M during sweep is independent of $N!$

④ For large D , use Krylov methods (Lanczos/Arnoldi) to minimize $A^{[s]}$!