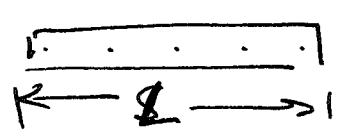


The area law

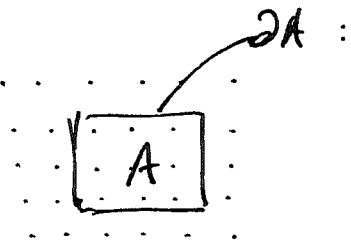
$$H = \sum h_i, \text{ gapped} \Rightarrow |4_0\rangle \text{ G.S.}$$

1D:

$|4_0\rangle \dots \overbrace{\dots}^{\beta_L} \dots \Rightarrow S(\beta_L) \leq \text{const.}$
 $(\text{proven - Hastings})$

even for gapless H typ. $S(\beta_L) \sim \log L!$
(but there exist examples w/ $S(\beta_L) \propto L!$)

2D:

 $S(S_A) \sim |\partial A|$
(or $|\partial A| \cdot \log |A|$)

→ scales like surface area vs. volume:

\Rightarrow Area law

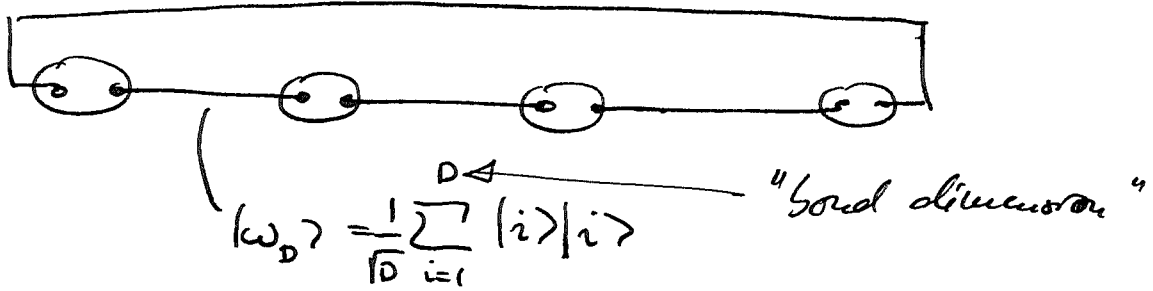
Intuition: Entanglement is concentrated around boundary (for any partition!)

Ansatz for area law (1D):

1. Decompose each site into 2 subsystems:



2. Put sites into max. entangled states ("bonds")

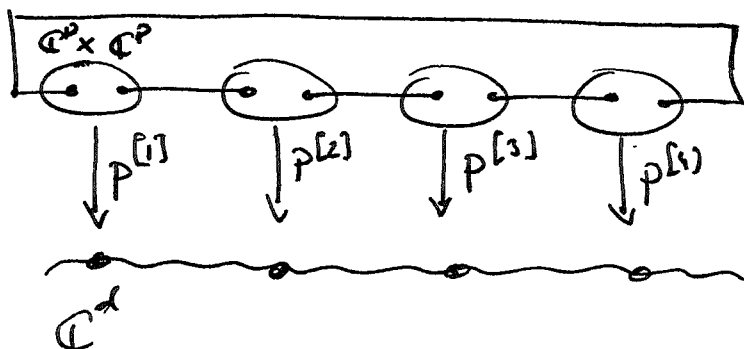


→ This ansatz has exactly an area law behavior:
for any partition, we cut exactly two bonds!

$$|S(\mathcal{S}_L) = 2 \log D|$$

~~2.~~ This is of course a bit silly: no NNN-correlations, ...

3. Apply a map $P^{[i]}$: $\mathbb{C}^D \times \mathbb{C}^D \rightarrow \mathbb{C}^d$ at each sites:



$$|\Psi\rangle = P^{[1]} \otimes \dots \otimes P^{[N]} |\omega_D\rangle^{\otimes N}$$

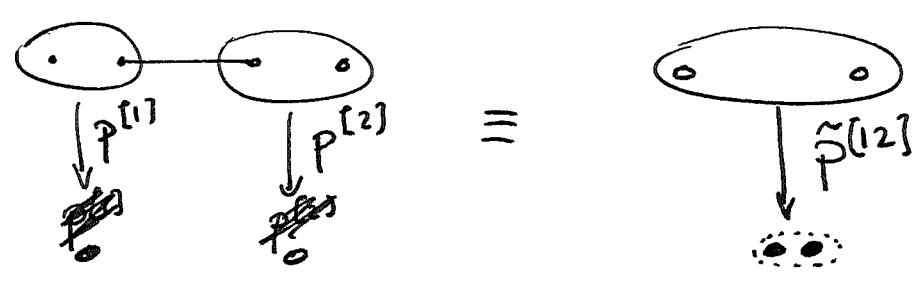
- This gives a family of states which can be enlarged by increasing D .
- $D=1$: Product states are special case.
- parameters: $N \cdot D^2 \cdot d$ \rightarrow efficient scaling in N .
- can approximate states w/ area law (and thus g.s. of gapped Hams.) with moderate scaling of D .

Matrix Product States:

What is the explicit form of $|\Psi\rangle$,

$$|\Psi\rangle = \sum c_{i_1, \dots, i_N} |i_1, \dots, i_N\rangle ?$$

Consider:



$$\equiv \left(\begin{matrix} P^{[1]} \\ \uparrow \\ A, B \end{matrix} \otimes \begin{matrix} P^{[2]} \\ \uparrow \\ C, D \end{matrix} \right) \cdot |\omega_0\rangle$$

\uparrow \uparrow \uparrow
 B, C

(18)

Write $P^{[s]} = \sum_{i\alpha\beta} A_{\alpha\beta}^{[s],i} |i\rangle\langle\alpha,\beta|$

$$\left(P_{AB}^{[1]} \otimes P_{CD}^{[2]} \right) |u_0\rangle_{BC} = \sum_{i\alpha\beta} A_{\alpha\beta}^{[1],i} |i\rangle\langle\alpha,\beta| \cdot \sum_{j\gamma\delta} A_{\gamma\delta}^{[2],j} |j\rangle\langle\gamma,\delta| \cdot \sum_{k=1}^D |k,k\rangle_{\beta,c}$$

$\beta = \gamma$

$$= \sum_{ij\alpha\beta\delta} A_{\alpha\beta}^{[1],i} A_{\beta\delta}^{[2],j} |ij\rangle\langle\alpha,\delta|$$

$$= \sum_{(ij)\alpha\delta} \left(A_{\alpha\beta}^{[1],i} A_{\beta\delta}^{[2],j} \right) |ij\rangle\langle\alpha,\delta|$$

$$=: \tilde{P}^{[12]} \leftarrow \text{Same form as before,}$$

with $A^{[s],i} \rightarrow A^{[s],i} A^{[s+1],i_{s+1}}$

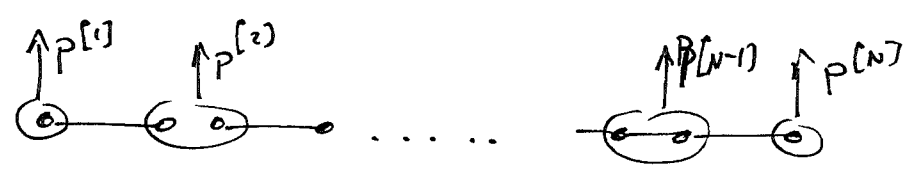
Iterate this argument:

Weight of config $|i_1, \dots, i_N\rangle$ is

$$e_{i_1 \dots i_N} = \text{tr} \left[A^{[1],i_1} A^{[2],i_2} \dots A^{[N],i_N} \right]$$

\Rightarrow Matrix Product States

* This can also be done with open boundary conditions (OBC)



$$|\psi\rangle = \sum_{i_1, \dots, i_N} \langle e | A^{[1], i_1} \cdot A^{[2], i_2} \cdot \dots \cdot A^{[N], i_N} | i_1, \dots, i_N \rangle$$

\uparrow \uparrow \uparrow
 $1 \times N$ $N \times N$ $N \times 1$

$$= \sum_{i_1, \dots, i_N} \langle e | \tilde{A}^{[1], i_1} \cdot A^{[2], i_2} \cdot \dots \cdot A^{[N], i_N} | e \rangle | i_1, \dots, i_N \rangle$$

* We can also extend this by allowing for different bond dimensions on each link:

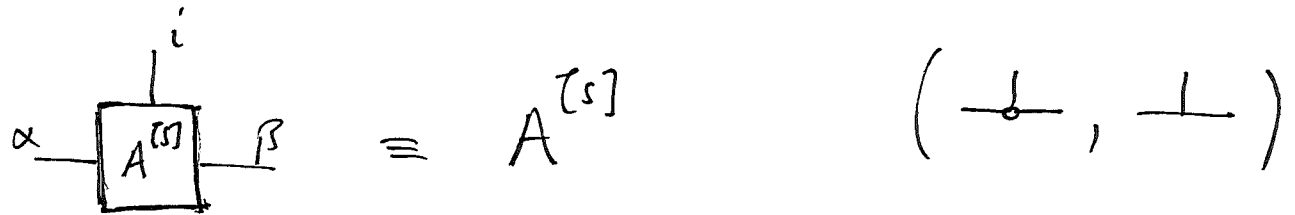
$$\sum_{i_1, \dots, i_N} \langle e | \left[A^{[1], i_1} \cdot A^{[2], i_2} \cdot \dots \cdot A^{[N], i_N} \right] | i_1, \dots, i_N \rangle$$

\uparrow \uparrow \uparrow
 $D_1 \times D_2$ $D_2 \times D_3$ $D_N \times D_1$

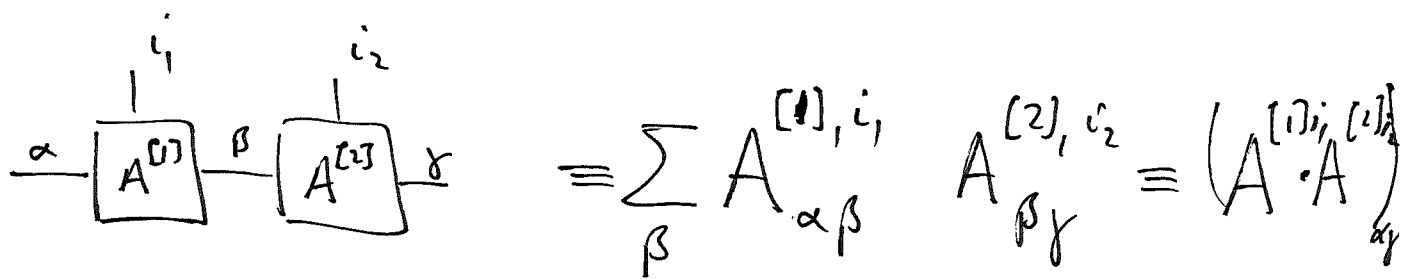
* Tensor network notation:

$$A^{[s]} \equiv A_{\alpha\beta}^{[s]i} : \text{tensor w/ 3 indices.}$$

Graphical notation:

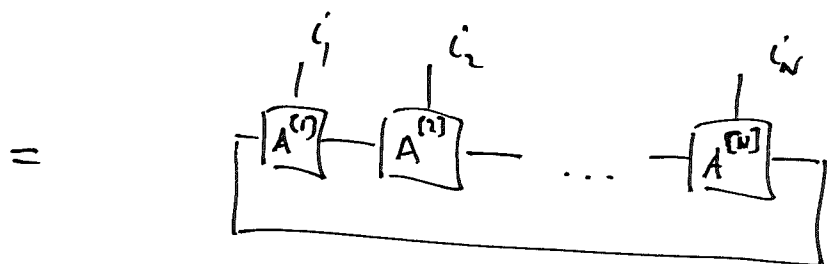


Summing over index ("contraction") \equiv connect legs:

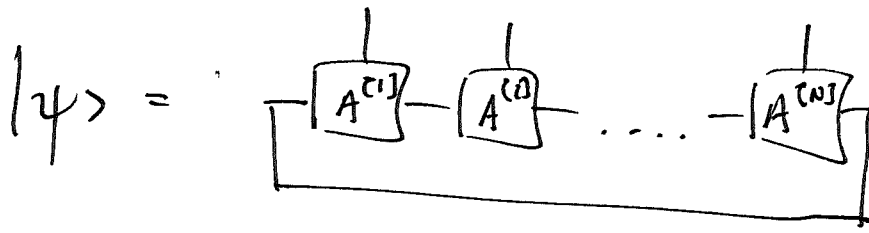


MPS: $|\psi\rangle = \sum c_{i_1 \dots i_N} |i_1, \dots, i_N\rangle$
 \uparrow
N-index tensor

$$c_{i_1 \dots i_N} = \text{tr} [A^{[1], i_1} \dots A^{[N], i_N}]$$



We will also write shorthand



Some examples:

* GHZ state: $|\psi\rangle = |0\dots 0\rangle + |1\dots 1\rangle$

- MPS (D=2): $A^0 = \begin{pmatrix} 1 & 0 \\ 0 & 0 \end{pmatrix}; A^1 = \begin{pmatrix} 0 & 0 \\ 0 & 1 \end{pmatrix};$ PBC
 (PBC)

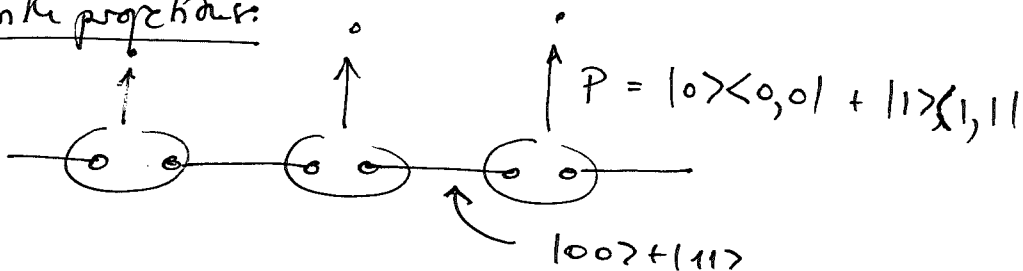
$A^0 A^0 = A^0$

$A^0 A^1 = 0$

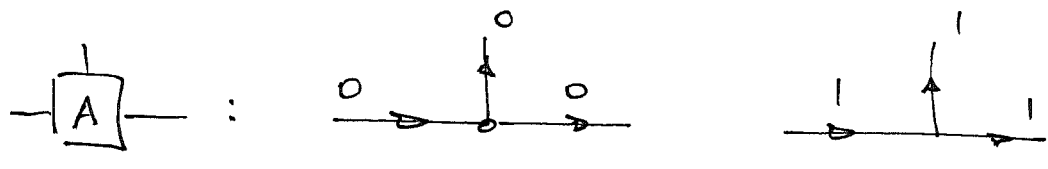
$A^1 A^1 = A^1$

$c_{i_1 \dots i_N} = \text{tr} [A^{i_1} \dots A^{i_N}] = \begin{cases} \text{tr}(A^0) = 1; i_1 = \dots = i_N = 0 \\ \text{tr}(A^1) = 1; i_1 = \dots = i_N = 1 \\ 0 \text{ otherwise} \end{cases}$

- with projectors:



"classical agents" picture:



(interpret tensors as "agents" passing information from left to right, and keeping some local info)

alternatively: OBC:

$$\langle l | = \langle + | ; \quad | r \rangle = | + \rangle$$

$$\langle + | A^1 \dots A^N | + \rangle = \begin{cases} \frac{1}{2}, & i_1 = i_2 = \dots = i_N \\ 0 & \text{otherwise} \end{cases}$$

* W state:

$$| \psi \rangle = | 10 \dots 0 \rangle + | 010 \dots \rangle + | 0010 \dots \rangle + \dots$$

NRPS ~~NRPS~~: $A^0 = \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix} = | 0 \rangle \langle 0 | + | 1 \rangle \langle 1 |$; $A^1 = \begin{pmatrix} 0 & 1 \\ 0 & 0 \end{pmatrix} = | 0 \rangle \langle 1 |$

D=2, OBC

$$\langle l | = \langle 0 | ; \quad | r \rangle = | 1 \rangle$$

Note: $A^0 \cdot A^0 = A^0$

$$A^0 \cdot A^1 = A^1 = |0 \times 1|$$

$$A^1 \cdot A^1 = 0$$

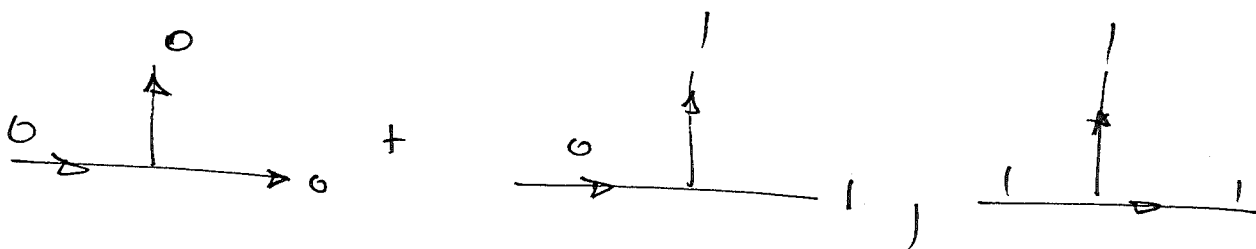
\Rightarrow at most one A^1 in string

$$A^{i_1} \dots A^{i_N} !$$

$$A^{i_1} \dots A^{i_N} = \begin{cases} A^0 & ; \text{ if } i_1 = \dots = i_N = 0 \\ A^1 & ; \text{ if } i_1 + i_2 + \dots + i_N = 1 \\ 0 & \text{ otherwise} \end{cases}$$

$$\Rightarrow \langle 0 | A^{i_1} \dots A^{i_N} | 1 \rangle = \begin{cases} \langle 0 | A^1 | 1 \rangle = 1 & ; \sum i_k = 1 \\ 0 & \text{ otherwise} \end{cases}$$

"agent" picture:



(pass bit from left to right & output 00, or output 1 and increase bit by one)

— bond space can be understood as counting the # of 1's which were output.

Note: No FBC rep. of W with fixed D leaves!