# Exercise Sheet 4

#### $Quantum \ Information$

# To be returned no later than May 15, 2015

### (20 points) Problem 1: Majorization.

1) Show that  $x \prec y$  if and only if for all real t,  $\sum_{j=1}^{d} \max\{x_j - t, 0\} \leq \sum_{j=1}^{d} \max\{y_j - t, 0\}$ , and  $\sum_{j=1}^{d} x_j = \sum_{j=1}^{d} y_j$ .

2) Use the previous problem to show that the set of x such that  $x \prec y$  is convex.

3) Give an example of two vectors x and y that are not related by majorization.

4) Let f be a convex function, and define  $F(x) = \sum f(x)$ . Show that  $x \prec y$  implies that  $F(x) \leq F(y)$  (a property known as "Schur convexity").

5) Use the previous problem to show that the entanglement  $E(|\psi\rangle) = S(\text{tr}_A|\psi\rangle\langle\psi|)$  cannot be increased by LOCC.

#### (25 points) **Problem 2: Equivalence of definitions.**

1) Prove that  $x \prec y$  if and only if  $x = \sum_j p_j P_j y$  for some probability distribution  $p_j$  and permutation matrices  $P_j$ .

<u>Hint</u>: Suppose that  $x \prec y$ . Use induction on the dimension d and suppose that x and y are d + 1dimensional vectors. The goal is to find a combination of permutations D such that a d-dimensional vector y' created by eliminating a first entry from the vector Dy majorizes vector  $x' = (x_2, \ldots, x_{d+1})$ . -Show that there exist j and  $T \in [0, 1]$  such that  $x_1 = ty_1 + (1 - t)y_j$ .

-Define D = tI + (1-t)T, where T is the permutation matrix which transposes the 1st and j-th matrix elements. Write the components of vector Dy.

-Define x' and y' by eliminating the first entry from x and Dy respectively. Show that  $x' \prec y'$ .

-Invoke the inductive hypothesis and finish the proof in this direction.

To prove the other direction, you may use Problem 1 above.

2) Show that  $x \prec y$  if and only if x = Dy for some doubly stochastic D.

<u>Hint</u>: You may use Birkhoff's theorem without proof, which states that any doubly stochastic D can be written as  $D = \sum p_j P_j$  with  $p_j$  a probability distribution and  $P_j$  permutations.

#### (20 points) Problem 3: LOCC transformations.

Suppose  $|\psi\rangle$  can be transformed to  $|\phi\rangle$  by LOCC. A general LOCC protocol can involve an arbitrary number of rounds of measurement and classical communication. In this problem, we will show that any LOCC protocol can be realized in a single round with only one-way communication, i.e., a protocol involving just the following steps: Alice performs a single measurement described by measurement operators  $K_j$ , sends the result j to Bob, who performs a unitary operation  $U_j$  on his system. 1) First, suppose that Bob performs a measurement with operators  $M_j = \sum_{kl} M_{j,kl} |k\rangle_B \langle l|_B$  on a pure state  $|\psi\rangle_{AB} = \sum \lambda_l |l\rangle_A |l\rangle_B$ , with resulting state denoted as  $|\psi_j\rangle$ . Now suppose that Alice performs a measurement with operators  $N_j = \sum_{kl} M_{j,kl} |k\rangle_A \langle l|_A$  on a pure state  $|\psi\rangle$ , with resulting state denoted as  $|\phi_j\rangle$ . Show that there exist unitaries  $U_j$  on system A and  $V_j$  on system B such that  $|\psi_j\rangle = (U_j \otimes V_j) |\phi_j\rangle$ .

2) Explain how a multi-round protocol can now be done with one measurement done by Alice and one unitary operation done by Bob conditioned on Alice's outcome.

## (20 points) Problem 4: Entanglement catalysis.

Suppose Alice and Bob share a pair of four level systems in the state  $|\psi\rangle = \sqrt{0.4} |00\rangle + \sqrt{0.4} |11\rangle + \sqrt{0.1} |22\rangle + \sqrt{0.1} |33\rangle$ . Show that it is not possible for them to convert this state by LOCC to the state  $|\phi\rangle = \sqrt{0.5} |00\rangle + \sqrt{0.25} |11\rangle + \sqrt{0.25} |22\rangle$ . Imagine, however, that a friendly bank is willing to offer them the loan of a catalyst, an entangled pair of qubits in the state  $|c\rangle = \sqrt{0.6} |00\rangle + \sqrt{0.4} |11\rangle$ . Show that it is possible for Alice and Bob to convert the state  $|\psi\rangle \otimes |c\rangle$  to  $|\phi\rangle \otimes |c\rangle$  by local operations and classical communication, allowing them to return the catalyst  $|c\rangle$  to the bank after the transformation is complete.

## (15 points) **Problem 5: Fidelity.**

1) Let a be a hermitian operator X, and let  $\lambda_i$  denote its eigenvalues. Then, the trace norm  $||X||_1 \equiv ||X||_{\text{tr}}$  is defined as  $||X||_1 := \sum_i |\lambda_i|$ , and the operator norm  $||X|| \equiv ||X||_{\infty}$  as  $||X||_{\infty} = \max_i |\lambda_i|$  (cf. Sheet 3, Problem 1). Show that for A, B hermitian,

$$|\operatorname{tr}(AB)| \le ||A||_1 || \, ||B||_{\infty}$$

2) Use the above inequality to derive the bound

$$\left|\left\langle\psi\right|O\left|\psi\right\rangle-\left\langle\phi\right|O\left|\phi\right\rangle\right|\leq 2\sqrt{1-|\langle\psi\left|\phi\right\rangle|^2}\;\|O\|_{\infty}\;.$$

(*Note:* The above inequality still holds for general operators, where the trace norm is defined as  $||X||_1 = \text{tr}|X|$ , where  $|X| := \sqrt{X^{\dagger}X}$ , and the operator norm as  $||X||_{\infty} := \lambda_{\max}(|X|)$ . It is a special case of the corresponding Hölder inequality  $|\text{tr}(AB^{\dagger})| \leq ||A||_p ||B||_q$  with 1/p + 1/q = 1, where  $||X||_p := (\text{tr}|X|^p)^{1/p}$ .)